

Slow and Stale Gradients Can Win the Race: Error-Runtime Trade-offs in Distributed SGD

Supplement

Sanghamitra Dutta
Carnegie Mellon
University

Gauri Joshi
Carnegie Mellon
University

Soumyadip Ghosh
IBM TJ Watson
Research Center

Parijat Dube
IBM TJ Watson
Research Center

Priya Nagpurkar
IBM TJ Watson
Research Center

6 STRONG CONVEXITY DISCUSSION

Definition 4 (Strong-Convexity). *A function $h(\mathbf{u})$ is defined to be c -strongly convex, if the following holds for all \mathbf{u}_1 and \mathbf{u}_2 in the domain:*

$$h(\mathbf{u}_2) \geq h(\mathbf{u}_1) + [\nabla h(\mathbf{u}_1)]^T (\mathbf{u}_2 - \mathbf{u}_1) + \frac{c}{2} \|\mathbf{u}_2 - \mathbf{u}_1\|_2^2.$$

For strongly convex functions, the following result holds for all \mathbf{u} in the domain of $h(\cdot)$.

$$2c(h(\mathbf{u}) - h^*) \leq \|\nabla h(\mathbf{u})\|_2^2. \quad (19)$$

The proof is derived in [Bottou et al., 2016]. For completeness, we give the sketch here.

Proof. Given a particular \mathbf{u} , let us define the quadratic function as follows:

$$q(\mathbf{u}') = h(\mathbf{u}) + \nabla h(\mathbf{u})^T (\mathbf{u}' - \mathbf{u}) + \frac{c}{2} \|\mathbf{u}' - \mathbf{u}\|_2^2$$

Now, $q(\mathbf{u}')$ is minimized at $\mathbf{u}' = \mathbf{u} - \frac{1}{c} \nabla h(\mathbf{u})$ and the value is $h(\mathbf{u}) - \frac{1}{2c} \|\nabla h(\mathbf{u})\|_2^2$. Thus, from the definition of strong convexity we now have,

$$\begin{aligned} h^* &\geq h(\mathbf{u}) + \nabla h(\mathbf{u})^T (\mathbf{u}' - \mathbf{u}) + \frac{c}{2} \|\mathbf{u}' - \mathbf{u}\|_2^2 \\ &\geq h(\mathbf{u}) - \frac{1}{2c} \|\nabla h(\mathbf{u})\|_2^2 \quad [\text{minimum value of } q(\mathbf{u}')]. \end{aligned}$$

□

7 RUNTIME ANALYSIS PROOFS

Here we provide all the proofs and supplementary information for all the results in Section 4.

7.1 Runtime of K -sync SGD

Proof of Lemma 3. We assume that the P learners have an i.i.d. computation times. When all the learners

start together, and we wait for the first K out of P i.i.d. random variables to finish, the expected computation time for that iteration is $\mathbb{E}[X_{K:P}]$, where $X_{K:P}$ denotes the K -th statistic of P i.i.d. random variables X_1, X_2, \dots, X_P . Thus, for J iterations, the runtime is given by $J\mathbb{E}[X_{K:P}]$. □

7.1.1 K -th statistic of exponential distributions

Here we give a sketch of why the K -th order statistic of P exponentials scales as $\log(P/P - K)$. A detailed derivation can be obtained in [Sheldon, 2002]. Consider P i.i.d. exponential distributions with parameter μ . The minimum $X_{1:P}$ of P independent exponential random variables with parameter μ is exponential with parameter $P\mu$. Conditional on $X_{1:P}$, the second smallest value $X_{2:P}$ is distributed like the sum of $X_{1:P}$ and an independent exponential random variable with parameter $(P - 1)\mu$. And so on, until the K -th smallest value $X_{K:P}$ which is distributed like the sum of $X_{(K-1):P}$ and an independent exponential random variable with parameter $(P - K + 1)\mu$. Thus,

$$X_{K:P} = Y_P + Y_{P-1} + \dots + Y_{P-K+1}$$

where the random variables Y_i s are independent and exponential with parameter $i\mu$. Thus,

$$\mathbb{E}[X_{K:P}] = \sum_{i=P-K+1}^P \frac{1}{i\mu} = \frac{H_P - H_{P-K}}{\mu} \approx \frac{\log \frac{P}{P-K}}{\mu}.$$

Here H_P and H_{P-K} denote the P -th and $(P - K)$ -th harmonic numbers respectively.

For the case where $K = P$, the expectation is given by,

$$\mathbb{E}[X_{P:P}] = \frac{1}{\mu} \sum_{i=1}^P \frac{1}{i} = \frac{1}{\mu} H_P \approx \frac{1}{\mu} \log P.$$

7.2 Runtime of K -batch-async SGD

Here we include a discussion on renewal processes for completeness, as a background to the proof of Lemma 4,

which gives the runtime of K -batch-async SGD. The familiar reader can skim through this part, and directly proceed to the proof of Lemma 4 in the main paper in Section 4.

Definition 5 (Renewal Process). *A renewal process is an arrival process where the inter-arrival intervals are positive, independent and identically distributed random variables.*

Lemma 6 (Elementary Renewal Theorem). [Gallager, 2013, Chapter 5] *Let $\{N(t), t > 0\}$ be a renewal counting process denoting the number of renewals in time t . Let $\mathbb{E}[Z]$ be the mean inter-arrival time. Then,*

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[N(t)]}{t} = \frac{1}{\mathbb{E}[Z]} \quad (20)$$

Observe that for asynchronous SGD or K -batch-async SGD, every gradient push by a learner to the PS can be thought of as an arrival process. The time between two consecutive pushes by a learner follows the distribution of X_i and is independent as computation time has been assumed to be independent across learners and mini-batches. Thus the inter-arrival intervals are positive, independent and identically distributed and hence, the gradient pushes are a renewal process.

7.3 Runtime of K -async SGD

Proof of Lemma 5. For new-longer-than-used distributions observe that the following holds:

$$\Pr(X_i > u + t | X_i > t) \leq \Pr(X_i > u) \quad (21)$$

Thus the random variable $X_i - t | X_i > t$ is thus stochastically dominated by X_i . Now let us assume we want to compute the expected computation time of one iteration of K -async starting at time instant t_0 . Let us also assume that the learners last read their parameter values at time instants t_1, t_2, \dots, t_P respectively where any K of these t_1, t_2, \dots, t_P are equal to t_0 as K out of P learners were updated at time t_0 and the remaining $(P - K)$ of these t_1, t_2, \dots, t_P are $< t_0$. Let Y_1, Y_2, \dots, Y_P be the random variables denoting the computation time of the P learners starting from time t_0 . Thus,

$$Y_i = X_i - (t_0 - t_i) | X_i > (t_0 - t_i) \quad \forall i = 1, 2, \dots, P \quad (22)$$

Now each of the Y_i s are independent and are stochastically dominated by the corresponding X_i s.

$$\Pr(Y_i > u) \leq \Pr(X_i > u) \quad \forall i, j = 1, 2, \dots, P \quad (23)$$

The expectation of the K -th statistic of $\{Y_1, Y_2, \dots, Y_P\}$ is the runtime of the iteration. Let us denote $h_K(x_1, x_2, \dots, x_P)$ as the K -th statistic of P numbers (x_1, x_2, \dots, x_P) . And let us denote

$g_{K,s}(x)$ as the K -th statistic of P numbers where $P - 1$ of them are given as $\mathbf{s}_{1 \times (P-1)}$ and x is the P -th number. Thus

$$g_{K,s}(x) = h_K(x, s(1), s(2), \dots, s(P-1)).$$

First observe that $g_{K,s}(x)$ is an increasing function of x since given the other $P - 1$ values, the K -th order statistic will either stay the same or increase with x . Now we use the property that if Y_i is stochastically dominated by X_i , then for any increasing function $g(\cdot)$, we have

$$\mathbb{E}_{Y_1} [g(Y_1)] \leq \mathbb{E}_{X_1} [g(X_1)].$$

This result is derived in [Kreps, 1990].

This implies that for a given \mathbf{s} ,

$$\mathbb{E}_{Y_1} [g_{K,s}(Y_1)] \leq \mathbb{E}_{X_1} [g_{K,s}(X_1)].$$

This leads to,

$$\begin{aligned} & \mathbb{E}_{Y_1 | Y_2=s(1), Y_3=s(2), \dots, Y_P=s(P-1)} [h_K(Y_1, Y_2, \dots, Y_P)] \\ & \leq \mathbb{E}_{X_1 | Y_2=s(1), Y_3=s(2), \dots, Y_P=s(P-1)} [h_K(X_1, Y_2, \dots, Y_P)] \end{aligned} \quad (24)$$

From this,

$$\begin{aligned} & \mathbb{E} [h_K(Y_1, Y_2, \dots, Y_P)] \\ & = \mathbb{E}_{Y_2, \dots, Y_P} [\mathbb{E}_{Y_1 | Y_2, Y_3, \dots, Y_P} [h_K(Y_1, Y_2, \dots, Y_P)]] \\ & \leq \mathbb{E}_{Y_2, \dots, Y_P} [\mathbb{E}_{X_1 | Y_2, Y_3, \dots, Y_P} [h_K(X_1, Y_2, \dots, Y_P)]] \\ & = \mathbb{E} [h_K(X_1, Y_2, \dots, Y_P)] \end{aligned} \quad (25)$$

This step proceeds inductively. Thus, similarly

$$\begin{aligned} & \mathbb{E} [h_K(X_1, Y_2, \dots, Y_P)] \\ & = \mathbb{E}_{X_1, Y_3, \dots, Y_P} [\mathbb{E}_{Y_2 | X_1, Y_3, \dots, Y_P} [h_K(X_1, Y_2, \dots, Y_P)]] \\ & \leq \mathbb{E}_{X_1, Y_3, \dots, Y_P} [\mathbb{E}_{X_2 | X_1, Y_3, \dots, Y_P} [h_K(X_1, X_2, Y_3, \dots, Y_P)]] \\ & = \mathbb{E} [h_K(X_1, X_2, Y_3, \dots, Y_P)] \end{aligned} \quad (26)$$

Thus, finally combining, we have,

$$\begin{aligned} & \mathbb{E} [h_K(Y_1, Y_2, \dots, Y_P)] \\ & \leq \mathbb{E} [h_K(X_1, Y_2, \dots, Y_P)] \\ & \leq \mathbb{E} [h_K(X_1, X_2, Y_3, \dots, Y_P)] \leq \dots \\ & \leq \mathbb{E} [h_K(X_1, X_2, X_3, \dots, X_P)] \end{aligned} \quad (27)$$

□

7.3.1 Special Case: Exponential Distributions

For exponential distributions, the inequality in Lemma 5 holds with equality. This follows from the memoryless property of exponentials. Let us consider

the scenario of the proof of Lemma 5 where we similarly define $Y_i = X_i - (t_0 - t_i) | X_i > (t_0 - t_i)$. From the memoryless property of exponentials [Sheldon, 2002], if $X_i \sim \exp(\mu)$, then $Y_i \sim \exp(\mu)$. Thus, the expectation of the K -th statistic of Y_i s can be easily derived as all the Y_i s are now i.i.d. with distribution $\exp(\mu)$. Thus, the runtime for J iterations is given by,

$$\mathbb{E}[T] = J\mathbb{E}[Y_{K:P}] = \frac{J}{\mu} \sum_{i=P-K+1}^P \frac{1}{i} \approx \frac{J}{\mu} \log \frac{P}{P-K}.$$

7.3.2 Comparison of K -async and K -batch-async SGD

We compare the error-runtime trade-off of K -async with K -batch-async SGD in Figure 10 as follows.

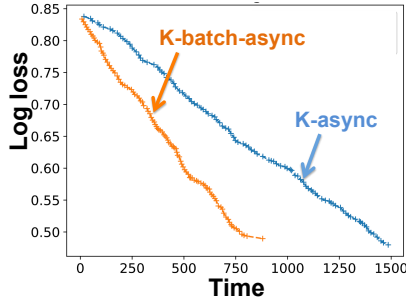


Figure 10: Accuracy Runtime Trade-off on MNIST Dataset: Comparison of K -async with K -batch-async under exponential computation time with $X_i \sim \exp(1)$. As derived theoretically, the K -batch-async has a sharper fall with time as compared to K -async even though the error attained is similar.

8 ASYNC-SGD ANALYSIS PROOFS

In this section, we provide a proof of the error convergence of asynchronous and K -async SGD.

8.1 Async-SGD with Fixed learning rate

First we prove a simplified version of Theorem 3 for the case $K = 1$. While this is actually a corollary of the more general Theorem 3, we prove this first for ease of understanding and simplicity. The proof of the more general Theorem 3 is then provided in Section 8.2.

Corollary 2. *Suppose that the objective function $F(\mathbf{w})$ is strongly convex with parameter c and the learning rate $\eta \leq \frac{1}{2L(\frac{M_G}{m} + 1)}$. Also assume that $\mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \leq \gamma \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2]$ for some constant $\gamma \leq 1$. Then, the error after J iterations*

of Async SGD is given by,

$$\mathbb{E}[F(\mathbf{w}_J)] - F^* \leq \frac{\eta L \sigma^2}{2c\gamma'm} + (1 - \eta c\gamma')^J (\mathbb{E}[F(\mathbf{w}_0)] - F^* - \frac{\eta L \sigma^2}{2c\gamma'm})$$

where $\gamma' = 1 - \gamma + \frac{p_0}{2}$ and p_0 is a non-negative lower bound on the conditional probability that $\tau(j) = j$ given all the past delays and parameters.

To prove the result, we will use the following lemma.

Lemma 7. *Let us denote $\mathbf{v}_j = g(\mathbf{w}_{\tau(j)}, \xi_j)$, and assume that $\mathbb{E}_{\xi_j|\mathbf{w}}[g(\mathbf{w}, \xi_j)] = \nabla F(\mathbf{w})$. Then,*

$$\mathbb{E}[\|\nabla F(\mathbf{w}_j) - \mathbf{v}_j\|_2^2] \leq \mathbb{E}[\|\mathbf{v}_j\|_2^2] - \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] + \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2].$$

Proof of Lemma 7. Observe that,

$$\begin{aligned} & \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \mathbf{v}_j\|_2^2] \\ &= \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)}) + \nabla F(\mathbf{w}_{\tau(j)}) - \mathbf{v}_j\|_2^2] \\ &= \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ & \quad + \mathbb{E}[\|\mathbf{v}_j - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \end{aligned} \quad (28)$$

The last line holds since the cross term is 0 as derived below.

$$\begin{aligned} & \mathbb{E}[(\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)}))^T (\mathbf{v}_j - \nabla F(\mathbf{w}_{\tau(j)}))] \\ &= \mathbb{E}_{\mathbf{w}_{\tau(j)}, \mathbf{w}_j} [(\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)}))^T \\ & \quad \mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}, \mathbf{w}_j} [(\mathbf{v}_j - \nabla F(\mathbf{w}_{\tau(j)}))]] \\ &= \mathbb{E}_{\mathbf{w}_{\tau(j)}, \mathbf{w}_j} [(\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)}))^T \\ & \quad (\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\mathbf{v}_j] - \nabla F(\mathbf{w}_{\tau(j)}))] = 0 \end{aligned}$$

Here again the last line follows from Assumption 2 in Section 2 which states that

$$\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\mathbf{v}_j] = \nabla F(\mathbf{w}_{\tau(j)}).$$

Returning to (28), observe that the second term can be further decomposed as,

$$\begin{aligned} & \mathbb{E}[\|\mathbf{v}_j - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ &= \mathbb{E}_{\mathbf{w}_{\tau(j)}} \left[\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\|\mathbf{v}_j - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \right] \\ &= \mathbb{E}_{\mathbf{w}_{\tau(j)}} \left[\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\|\mathbf{v}_j\|_2^2] \right. \\ & \quad \left. - 2\mathbb{E}_{\mathbf{w}_{\tau(j)}} \left[\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\mathbf{v}_j^T \nabla F(\mathbf{w}_{\tau(j)})] \right] \right. \\ & \quad \left. + \mathbb{E}_{\mathbf{w}_{\tau(j)}} \left[\mathbb{E}_{\xi_j|\mathbf{w}_{\tau(j)}} [\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \right] \right] \\ &= \mathbb{E}[\|\mathbf{v}_j\|_2^2] - 2\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] + \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ &= \mathbb{E}[\|\mathbf{v}_j\|_2^2] - \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2]. \end{aligned}$$

□

We also prove a K -learner version of this lemma to prove Theorem 3. Now we proceed to provide the proof of Corollary 2.

Proof of Corollary 2.

$$\begin{aligned}
 F(\mathbf{w}_{j+1}) &\leq F(\mathbf{w}_j) + (\mathbf{w}_{j+1} - \mathbf{w}_j)^T \nabla F(\mathbf{w}_j) \\
 &\quad + \frac{L}{2} \|\mathbf{w}_{j+1} - \mathbf{w}_j\|_2^2 \\
 &= F(\mathbf{w}_j) + (-\eta \mathbf{v}_j)^T \nabla F(\mathbf{w}_j) + \frac{L\eta^2}{2} \|\mathbf{v}_j\|_2^2 \\
 &= F(\mathbf{w}_j) - \frac{\eta}{2} \|\nabla F(\mathbf{w}_j)\|_2^2 - \frac{\eta}{2} \|\mathbf{v}_j\|_2^2 \\
 &\quad + \frac{\eta}{2} \|\nabla F(\mathbf{w}_j) - \mathbf{v}_j\|_2^2 + \frac{L\eta^2}{2} \|\mathbf{v}_j\|_2^2 \quad (29)
 \end{aligned}$$

Here the last line follows from $2\mathbf{a}^T \mathbf{b} = \|\mathbf{a}\|_2^2 + \|\mathbf{b}\|_2^2 - \|\mathbf{a} - \mathbf{b}\|_2^2$. Taking expectation,

$$\begin{aligned}
 \mathbb{E}[F(\mathbf{w}_{j+1})] &\leq \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\
 &\quad - \frac{\eta}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] + \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \mathbf{v}_j\|_2^2] \\
 &\quad + \frac{L\eta^2}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\
 &\stackrel{(a)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] - \frac{\eta}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\
 &\quad + \frac{\eta}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] - \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\
 &\quad + \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\
 &\quad + \frac{L\eta^2}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \quad (30)
 \end{aligned}$$

Here, (a) follows from Lemma 7 that we just derived. Now, again bounding from (30), we have

$$\begin{aligned}
 \mathbb{E}[F(\mathbf{w}_{j+1})] &\quad (31) \\
 &\stackrel{(b)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] - \frac{\eta}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\
 &\quad + \frac{\eta}{2} \gamma \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\
 &\stackrel{(c)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} (1 - \gamma) \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2 \sigma^2}{2m} \\
 &\quad - \frac{\eta}{2} \left(1 - L\eta \left(\frac{M_G}{m} + 1\right)\right) \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\
 &\stackrel{(d)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} (1 - \gamma) \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2 \sigma^2}{2m} \\
 &\quad - \frac{\eta}{4} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\
 &\stackrel{(e)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2} (1 - \gamma) \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2 \sigma^2}{2m} \\
 &\quad - \frac{\eta}{4} p_0 \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \quad (32)
 \end{aligned}$$

Here (b) follows from the statement of the theorem that

$$\mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \leq \gamma \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2]$$

for some constant $\gamma \leq 1$. The next step (c) follows from Assumption 4 in Section 2 which lead to

$$\mathbb{E}[\|\mathbf{v}_j\|_2^2] \leq \frac{\sigma^2}{m} + \left(\frac{M_G}{m} + 1\right) \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2].$$

Step (d) follows from choosing $\eta < \frac{1}{2L(\frac{M_G}{m} + 1)}$ and finally (e) follows from Lemma 1.

Now one might recall that the function $F(w)$ was defined to be strongly convex with parameter c . Using the standard result of strong-convexity (6) in (32), we obtain the following result.

$$\begin{aligned}
 \mathbb{E}[F(\mathbf{w}_{j+1})] - F^* &\leq \frac{\eta^2 L \sigma^2}{2m} \\
 &\quad + (1 - \eta c (1 - \gamma + \frac{p_0}{2})) (\mathbb{E}[F(\mathbf{w}_j)] - F^*)
 \end{aligned}$$

Let us denote $\gamma' = (1 - \gamma + \frac{p_0}{2})$. Then, using the above recursion, we thus have,

$$\begin{aligned}
 \mathbb{E}[F(\mathbf{w}_J)] - F^* &\leq \frac{\eta L \sigma^2}{2c\gamma' m} + \\
 &\quad (1 - \eta\gamma' c)^J (\mathbb{E}[F(\mathbf{w}_0)] - F^* - \frac{\eta L \sigma^2}{2c\gamma' m})
 \end{aligned}$$

□

8.1.1 Discussion on range of p_0

Let us denote the conditional probability of $\tau(j) = j$ given all the past delays and parameters as $p_0^{(j)}$. Now $p_0 \leq p_0^{(j)} \forall j$. Clearly the value of $p_0^{(j)}$ will differ for different distributions and accordingly the value of p_0 will differ. Here we include a brief discussion on the possible values of p_0 for different distributions. These also hold for K -async and K -batch-async SGD.

Lemma 8 (Bounds of p_0). *Define $p_0 = \inf_j p_0^{(j)}$, i.e. the largest constant such that $p_0 \leq p_0^{(j)} \forall j$.*

- For exponential computation times, $p_0^{(j)} = \frac{1}{P}$ for all j and is thus invariant of j and $p_0 = \frac{1}{P}$.
- For new-longer-than-used (See Definition 3) computation times, $p_0^{(j)} \leq \frac{1}{P}$ and thus $p_0 \leq \frac{1}{P}$.
- For new-shorter-than-used computation times, $p_0^{(j)} \geq \frac{1}{P}$ and thus $p_0 \geq \frac{1}{P}$.

Proof of Lemma 8. Let t_0 be the time when the j -th iteration occurs, and suppose that learner i' pushed its gradient in the j -th iteration. Now similar to the

proof of Lemma 5, let us also assume that the learners last read their parameter values at time instants t_1, t_2, \dots, t_P respectively where $t'_i = t_0$ and the remaining $(P - 1)$ of these t_i s are $< t_0$. Let Y_1, Y_2, \dots, Y_P be the random variables denoting the computation time of the P learners starting from time t_0 . Thus, $Y_i = X_i - (t_0 - t_i) | X_i > (t_0 - t_i)$. For exponentials, from the memoryless property, all these Y_i s become i.i.d. and thus from symmetry the probability of i' finishing before all the others is equal, *i.e.* $\frac{1}{P}$. Thus, $p_0^{(j)} = p_0 = \frac{1}{P}$. For new-longer-than-used distributions, as we have discussed before all the Y_i s with $i \neq i'$ will be stochastically dominated by $Y_{i'} = X_{i'}$. Thus, probability of i s with $i \neq i'$ finishing first is higher than i' . Thus, $p_0^{(j)} \leq \frac{1}{P}$ and so is p_0 . Similarly, for new-shorter-than-used distributions, $Y_{i'}$ is stochastically dominated by all the Y_i s and thus probability of i' finishing first is more. So, $p_0^{(j)} \geq \frac{1}{P}$ and so is p_0 . \square

8.2 K-async SGD under fixed learning rate

In this subsection, we provide a proof of Theorem 3.

Before we proceed to the proof of this theorem, we first extend our Assumption 4 from the variance of a single stochastic gradient to sum of stochastic gradients in the following Lemma.

Lemma 9. *If the variance of the stochastic updates is bounded as $\mathbb{E}_{\xi_j | \mathbf{w}_{\tau(l,j)}} [\|g(\mathbf{w}_{\tau(l,j)}, \xi_{l,j}) - \nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \leq \frac{\sigma^2}{m} + \frac{M_G}{m} \|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2 \forall \tau(l,j) \leq j$, then for K -async, the variance of the sum of stochastic updates given all the parameter values $\mathbf{w}_{\tau(l,j)}$ is also bounded as follows:*

$$\begin{aligned} & \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \mathbf{w}_{\tau(1,j)} \dots \mathbf{w}_{\tau(K,j)}} \left[\left\| \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j}) \right\|_2^2 \right] \\ & \leq \frac{K\sigma^2}{m} + \left(\frac{M_G}{m} + K \right) \left\| \sum_{l=1}^K \nabla F(\mathbf{w}_{\tau(l,j)}) \right\|_2^2 \quad (33) \end{aligned}$$

Proof. First let us consider the expectation of any cross term such that $l \neq l'$. For the ease of writing, let $\Omega = \{\mathbf{w}_{\tau(1,j)} \dots \mathbf{w}_{\tau(K,j)}\}$. Now observe the conditional expectation of the cross term as follows.

$$\begin{aligned} & \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} [(g(\mathbf{w}_{l,j}, \xi_{l,j}) - \nabla F(\mathbf{w}_{\tau(l,j)}))^T \\ & \quad ((g(\mathbf{w}_{l',j}, \xi_{l',j}) - \nabla F(\mathbf{w}_{\tau(l',j)})))] \\ & = \mathbb{E}_{\xi_{l,j}, \xi_{l',j} | \Omega} [(g(\mathbf{w}_{l,j}, \xi_{l,j}) - \nabla F(\mathbf{w}_{\tau(l,j)}))^T \\ & \quad ((g(\mathbf{w}_{l',j}, \xi_{l',j}) - \nabla F(\mathbf{w}_{\tau(l',j)})))] \\ & = \mathbb{E}_{\xi_{l',j} | \Omega} [\mathbb{E}_{\xi_{l,j} | \xi_{l',j}, \Omega} [(g(\mathbf{w}_{l,j}, \xi_{l,j}) - \nabla F(\mathbf{w}_{\tau(l,j)}))^T \\ & \quad (g(\mathbf{w}_{l',j}, \xi_{l',j}) - \nabla F(\mathbf{w}_{\tau(l',j)}))]] \\ & = \mathbb{E}_{\xi_{l',j} | \Omega} [0^T (g(\mathbf{w}_{l',j}, \xi_{l',j}) - \nabla F(\mathbf{w}_{\tau(l',j)}))] = 0 \quad (34) \end{aligned}$$

Thus the cross terms are all 0. So the expression

simplifies as,

$$\begin{aligned} & \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} \left[\left\| \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j}) - F(\mathbf{w}_{\tau(l,j)}) \right\|_2^2 \right] \\ & \stackrel{(a)}{=} \sum_{l=1}^K \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} [\|g(\mathbf{w}_{l,j}, \xi_{l,j}) - F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ & \leq \sum_{l=1}^K \frac{\sigma^2}{m} + \frac{M_G}{m} \|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2 \quad (35) \end{aligned}$$

Thus,

$$\begin{aligned} & \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} \left[\left\| \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j}) \right\|_2^2 \right] \\ & = \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} \left[\left\| \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j}) - F(\mathbf{w}_{\tau(l,j)}) \right\|_2^2 \right] \\ & \quad + \mathbb{E}_{\xi_{1,j}, \dots, \xi_{K,j} | \Omega} \left[\left\| \sum_{l=1}^K F(\mathbf{w}_{\tau(l,j)}) \right\|_2^2 \right] \\ & \leq \frac{K\sigma^2}{m} + \sum_{l=1}^K \frac{M_G}{m} \|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2 + \left\| \sum_{l=1}^K F(\mathbf{w}_{\tau(l,j)}) \right\|_2^2 \\ & \leq \frac{K\sigma^2}{m} + \sum_{l=1}^K \frac{M_G}{m} \|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2 + \sum_{l=1}^K K \|F(\mathbf{w}_{\tau(l,j)})\|_2^2 \quad (36) \end{aligned}$$

Now we return to the proof of the theorem. \square

Proof of Theorem 3. Let $\mathbf{v}_j = \frac{1}{K} \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j})$. Following steps similar to the Async-SGD proof, from Lipschitz continuity we have the following.

$$\begin{aligned} & F(\mathbf{w}_{j+1}) \leq F(\mathbf{w}_j) + (\mathbf{w}_{j+1} - \mathbf{w}_j)^T \nabla F(\mathbf{w}_j) \\ & \quad + \frac{L}{2} \|\mathbf{w}_{j+1} - \mathbf{w}_j\|_2^2 \\ & = F(\mathbf{w}_j) - \frac{\eta}{K} \sum_{l=1}^K g(\mathbf{w}_{l,j}, \xi_{l,j})^T \nabla F(\mathbf{w}_j) + \frac{L}{2} \|\eta \mathbf{v}_j\|_2^2 \\ & \stackrel{(a)}{=} F(\mathbf{w}_j) - \frac{\eta}{2K} \sum_{l=1}^K \|\nabla F(\mathbf{w}_j)\|_2^2 - \frac{\eta}{2K} \sum_{l=1}^K \|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2 \\ & \quad + \frac{\eta}{2K} \sum_{l=1}^K \|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2 - \frac{\eta}{2K} \sum_{l=1}^K \|\nabla F(\mathbf{w}_j)\|_2^2 \\ & \quad + \frac{L\eta^2}{2} \|\mathbf{v}_j\|_2^2 \\ & = F(\mathbf{w}_j) - \frac{\eta}{2} \|\nabla F(\mathbf{w}_j)\|_2^2 - \frac{\eta}{2K} \sum_{l=1}^K \|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2 \\ & \quad + \frac{\eta}{2K} \sum_{l=1}^K \|g(\mathbf{w}_{l,j}, \xi_{l,j}) - \nabla F(\mathbf{w}_j)\|_2^2 \\ & \quad + \frac{L\eta^2}{2} \|\mathbf{v}_j\|_2^2 \quad (37) \end{aligned}$$

Here (a) follows from $2\mathbf{a}^T\mathbf{b} = \|\mathbf{a}\|_2^2 + \|\mathbf{b}\|_2^2 - \|\mathbf{a} - \mathbf{b}\|_2^2$. Taking expectation,

$$\begin{aligned} \mathbb{E}[F(\mathbf{w}_{j+1})] &\leq \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad - \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2] \\ &\quad + \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|\nabla F(\mathbf{w}_j) - g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2] \\ &\quad + \frac{L\eta^2}{2}\mathbb{E}[\|\mathbf{v}_j\|_2^2] \\ &\stackrel{(a)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad - \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2] + \\ &\quad \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|g(\mathbf{w}_{l,j}, \xi_{l,j})\|_2^2] \\ &\quad - \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ &\quad + \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ &\quad + \frac{L\eta^2}{2}\mathbb{E}[\|\mathbf{v}_j\|_2^2] \quad (38) \end{aligned}$$

$$\begin{aligned} &\stackrel{(b)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad - \frac{\eta}{2K}\sum_{l=1}^K\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ &\quad + \frac{\eta}{2}\gamma\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2}{2}\mathbb{E}[\|\mathbf{v}_j\|_2^2] \\ &\stackrel{(c)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}(1-\gamma)\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2\sigma^2}{2Km} \\ &\quad - \frac{\eta}{2K}\sum_{l=1}^K\left(1 - L\eta\left(\frac{M_G}{Km} + \frac{1}{K}\right)\right)\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ &\stackrel{(d)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}(1-\gamma)\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2\sigma^2}{2Km} \\ &\quad - \frac{\eta}{4K}\sum_{l=1}^K\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \\ &\stackrel{(e)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta}{2}(1-\gamma)\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{L\eta^2\sigma^2}{2Km} \\ &\quad - \frac{\eta}{4}p_0\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \quad (39) \end{aligned}$$

Here step (a) follows from Lemma 7 and step (b) follows from the assumption that $\mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \leq \gamma\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2]$ for some constant $\gamma \leq 1$. The next step (c) follows from the Lemma 9 that bounds the variance of the sum of stochastic gradients. Step (d) fol-

lows from choosing $\eta < \frac{1}{2L(\frac{M_G}{Km} + \frac{1}{K})}$ and finally (e) follows from Lemma 1 in Section 3 that says $\mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(l,j)})\|_2^2] \geq p_0\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2]$ for some non-negative constant p_0 which is a lower bound on the conditional probability that $\tau(l, j) = j$ given all past delays and parameter values.

Finally, since $F(\mathbf{w})$ is strongly convex, using the inequality $2c(F(\mathbf{w}) - F^*) \leq \|\nabla F(\mathbf{w})\|_2^2$ in (39), we finally obtain the desired result. \square

8.2.1 Extension to Non-Convex case

The analysis can be extended to provide weaker guarantees for non-convex objectives. Let $\gamma' = 1 - \gamma + \frac{p_0}{2}$

For non-convex objectives, we have the following result.

Theorem 5. *For non-convex objective function, we have the following ergodic convergence result given by:*

$$\frac{1}{J+1}\sum_{j=0}^J\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \leq \frac{2(F(\mathbf{w}_0) - F^*)}{(J+1)\eta\gamma'} + \frac{L\eta\sigma^2}{Km\gamma'}$$

where $F^* = \min_{\mathbf{w}} F(\mathbf{w})$.

Proof. Recall the recursion derived in the last proof in (39). After re-arrangement, we obtain the following:

$$\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \leq \frac{2(\mathbb{E}[F(\mathbf{w}_j)] - \mathbb{E}[F(\mathbf{w}_{j+1})])}{\eta\gamma'} + \frac{L\eta\sigma^2}{Km\gamma'} \quad (40)$$

Taking summation from $j = 0$ to $j = J$, we get,

$$\begin{aligned} &\frac{1}{J+1}\sum_{j=0}^J\mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\leq \frac{2(\mathbb{E}[F(\mathbf{w}_0)] - \mathbb{E}[F(\mathbf{w}_J)])}{(J+1)\eta\gamma'} + \frac{L\eta\sigma^2}{Km\gamma'} \\ &\stackrel{(a)}{\leq} \frac{2(F(\mathbf{w}_0) - F^*)}{(J+1)\eta\gamma'} + \frac{L\eta\sigma^2}{Km\gamma'} \quad (41) \end{aligned}$$

Here (a) follows since we assume \mathbf{w}_0 to be known and also from $\mathbb{E}[F(\mathbf{w}_J)] \geq F^*$. \square

8.3 Variable Learning Rate Schedule

We propose a new heuristic for learning rate schedule that is more stable than fixed learning rate for asynchronous SGD. Our learning rate schedule is $\eta_j = \min\left\{\frac{C}{\|\mathbf{w}_j - \mathbf{w}_{\tau(j)}\|_2^2}, \eta_{max}\right\}$, where η_{max} is a suitably large value of learning rate beyond which the convergence diverges. This heuristic is inspired from the assumption in Theorem 4 given by $\eta_j\mathbb{E}[\|\mathbf{w}_j - \mathbf{w}_{\tau(j)}\|_2^2] \leq C$. In this section, we derive the accuracy trade-off mentioned in Theorem 4 based on this assumption.

Proof of Theorem 4. Following steps similar to (29), we first obtain the following:

$$F(\mathbf{w}_{j+1}) \leq F(\mathbf{w}_j) - \frac{\eta_j}{2} \|\nabla F(\mathbf{w}_j)\|_2^2 - \frac{\eta_j}{2} \|\mathbf{v}_j\|_2^2 + \frac{\eta_j}{2} \|\nabla F(\mathbf{w}_j) - \mathbf{v}_j\|_2^2 + \frac{L\eta_j^2}{2} \|\mathbf{v}_j\|_2^2 \quad (42)$$

Now taking expectation, we obtain the following result. $\mathbb{E}[F(\mathbf{w}_{j+1})]$

$$\begin{aligned} &\stackrel{(a)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] - \frac{\eta_j}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\ &\quad + \frac{\eta_j}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ &\quad + \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j) - \nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ &\quad + \frac{L\eta_j^2}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\ &\stackrel{(b)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] + \frac{CL^2}{2} + \frac{L\eta_j^2}{2} \mathbb{E}[\|\mathbf{v}_j\|_2^2] \\ &\stackrel{(c)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \frac{CL^2}{2} + \frac{L\eta_j^2\sigma^2}{2m} \\ &\quad - \frac{\eta_j}{2} \left(1 - L\eta_j \left(\frac{M_G}{m} + 1\right)\right) \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \\ &\stackrel{(e)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad + \frac{CL^2}{2} + \frac{\eta_j^2 L\sigma^2}{2m} - \frac{\eta_j}{4} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] \quad (43) \end{aligned}$$

Here (a) follows from (30), (b) follows from (12), (c) follows from Assumption 4 and (d) follows as $\eta_j \leq \frac{1}{2L(\frac{M_G}{m} + 1)}$. Let us define $\Delta_j = \frac{CL^2}{2} + \frac{\eta_j^2 L\sigma^2}{2m}$. Thus, the recursion can be written as,

$$\begin{aligned} \mathbb{E}[F(\mathbf{w}_{j+1})] &\leq \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] \\ &\quad - \frac{\eta_j}{4} \mathbb{E}[\|\nabla F(\mathbf{w}_{\tau(j)})\|_2^2] + \Delta_j \\ &\stackrel{(e)}{\leq} \mathbb{E}[F(\mathbf{w}_j)] - \frac{\eta_j}{2} \left(1 + \frac{p_0}{2}\right) \mathbb{E}[\|\nabla F(\mathbf{w}_j)\|_2^2] + \Delta_j \quad (44) \end{aligned}$$

Here (e) follows from Lemma 1. If the loss function $F(\mathbf{w})$ is strongly convex with parameter c , then for all \mathbf{w} , we have $2c(F(\mathbf{w}) - F^*) \leq \|\nabla F(\mathbf{w})\|_2^2$. Using this result, we obtain

$$\begin{aligned} \mathbb{E}[F(\mathbf{w}_{j+1})] - F^* &\leq (1 - \eta_j(1 + \frac{p_0}{2})c)(\mathbb{E}[F(\mathbf{w}_j)] - F^*) \\ &\quad + \Delta_j \\ &\leq (1 - \eta_j(1 + \frac{p_0}{2})c)(1 - \eta_{j-1}(1 + \frac{p_0}{2})c)(\mathbb{E}[F(\mathbf{w}_{j-1})] - F^*) \\ &\quad + (1 - \eta_j(1 + \frac{p_0}{2})c)\Delta_{j-1} + \Delta_j \\ &\leq (1 - \rho_j)(1 - \rho_{j-1}) \dots (1 - \rho_0)(\mathbb{E}[F(\mathbf{w}_0)] - F^*) + \Delta \quad (45) \end{aligned}$$

where $\rho_j = \eta_j(1 + \frac{p_0}{2})c$ and $\Delta = \Delta_j + (1 - \rho_j)\Delta_{j-1} + \dots + (1 - \rho_j)(1 - \rho_{j-1}) \dots (1 - \rho_1)\Delta_0$. \square

9 SIMULATION SETUP DETAILS

MNIST [LeCun, 1998]: For the simulations on MNIST dataset, we first convert the 28×28 images into single vectors of length 784. We use a single layer of neurons followed by soft-max cross entropy with logits loss function. Thus effectively the parameters consist of a weight matrix \mathbf{W} of size 784×10 and a bias vector \mathbf{b} of size 1×10 . We use a regularizer of value 0.01, mini-batch size $m = 1$, and learning rate $\eta = 0.01$. For implementation we used Tensorflow with Python3. Thus, the model is as follows:

```
X=tf.placeholder(tf.float32,[None,784])
Y=tf.placeholder(tf.float32,[None,10])
W=tf.Variable(tf.random_normal(shape=[784,10],
                                stddev=0.01), name="weights")
b=tf.Variable(tf.random_normal(shape=[1,10],
                                stddev=0.01), name="bias")

logits=tf.matmul(X,W) + b
entropy=tf.nn.softmax_cross_entropy_with
    _logits(logits=logits,labels=Y) +
    lamda*tf.square(tf.norm(W))

loss=tf.reduce_mean( entropy)
```

For the run-time simulations, we generate random variables from the respective distributions in python to represent the computation times.

CIFAR10 [Krizhevsky and Hinton, 2009]: For the CIFAR10 simulations, similar to MNIST, we convert the images into vectors of length 1024. We combine the three colour variants in the ratio [0.2989, 0.5870, 0.114] to generate a single vector of length 1024 for every image. We use a single layer of neurons again followed by soft-max cross entropy with logits in tensorflow. Thus, the parameters consist of a weight matrix \mathbf{W} of size 1024×10 and a bias vector \mathbf{b} of size 1×10 . We use a mini-batch size of 250, regularizer of 0.05.

We use a similar model as follows:

```
X=tf.placeholder(tf.float32,[None,1024])
Y=tf.placeholder(tf.float32,[None,10])
W=tf.Variable(tf.random_normal(shape=[1024,10],
                                stddev= 0.01),name="weights")
b=tf.Variable(tf.random_normal(shape=[1,10],
                                stddev = 0.01),name="bias")

logits=tf.matmul(X,W) + b
entropy=tf.nn.softmax_cross_entropy_with
    _logits(logits=logits,labels=Y) +
```

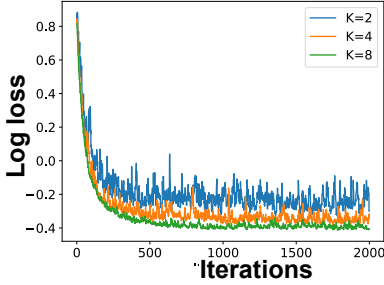


Figure 11: Error-Iterations tradeoff on MNIST dataset: Simulation of K -sync SGD for different values of K . Observe that accuracy improves with increasing K which means increasing effective batch size ($\eta = 0.05$).

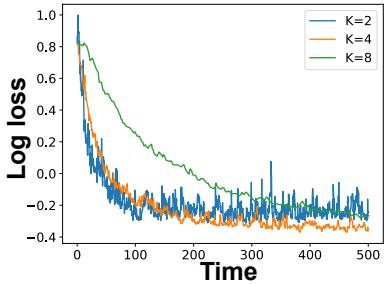


Figure 12: Error-Runtime tradeoff on MNIST dataset: Simulation of K -sync SGD for different values of K ($\eta = 0.05$).

```
lamda*tf.square(tf.norm(W))
loss=tf.reduce_mean(entropy)
```

The computation time as each learner is generated from exponential distribution.

10 CHOICE OF HYPERPARAMETERS

Our analysis techniques can also inform the choice of hyperparameters for synchronous and K -sync SGD.

10.1 Varying K in K -sync

We first perform some simulations of K -sync SGD applied on the MNIST dataset. For the simulation setup, we consider 8 parallel learners with fixed mini-batch size $m = 1$ and fixed learning rate 0.05. The number of learners to wait for in K -sync, *i.e.* K is varied and the error-runtime trade-off is observed. The runtimes are generated from a shifted exponential distribution given by $X_i \sim m + \exp \mu$.

Observe that in the plot of error with the number of iterations in Figure 11, the error improves with increasing K , which means increasing the effective mini-batch and reducing the variability in the gradient. However, if we look at the same error plotted against runtime

(See Figure 12) instead of the number of iterations, observe that increasing K naively does not always lead to a better trade-off. As K increases, the central PS has to wait for more learners to finish at every iteration, thus suffering from increased straggler effect. The best error-runtime trade-off is obtained at an intermediate $K = 4$. Thus, the current analysis informs the optimal choice of K to achieve a good error-runtime trade-off.

10.2 Varying mini-batch m

We consider the training of Alexnet on ImageNet dataset [Krizhevsky et al., 2012] using $P = 4$ learners. For this simulation, we perform fully synchronous SGD, *i.e.* K -sync with $K = P = 4$. We fix the learning rate and vary the mini-batch used for training. The runtimes are generated from a shifted exponential distribution given by $X_i \sim m + \exp \mu$, that depends on the mini-batch size. Intuitively, this distribution makes sense since to compute one mini-batch, a learner would atleast need a time m (Work Complexity). However, due to delays, it has the additional exponential tail. The error-runtime trade-offs are observed in Figure 13 and Figure 14.

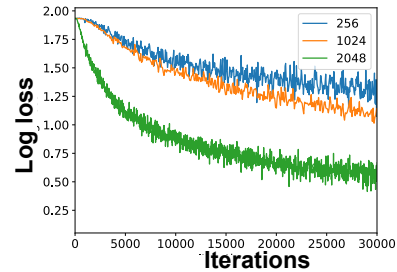


Figure 13: Error-Iterations tradeoff on IMAGENET dataset: Simulation of fully synchronous SGD ($K = P = 4$) for different values of mini-batch m . Observe that accuracy improves with increasing m which means increasing effective batch size.

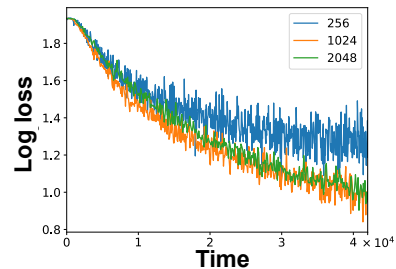


Figure 14: Error-Runtime tradeoff on IMAGENET dataset: Same simulation of fully synchronous SGD ($K = P = 4$) for different values of mini-batch m plotted against time. Observe that higher m does not necessarily mean the best trade-off with runtime as higher mini-batch also has longer time.

Again, observe that the plot of error with the number of iterations improves with the mini-batch size, as also expected from theory. However, increasing the mini-batch also changes the runtime distribution. Thus, when we plot the same error against runtime, we again observe that increasing the mini-batch size naively does not necessarily lead to the best trade-off. Instead, the best error-runtime trade-off is observed with an intermediate mini-batch value of 1024. Thus, our analysis informs the choice of the optimal mini-batch.